

A Foundation for Watermarking in Compressed Domain

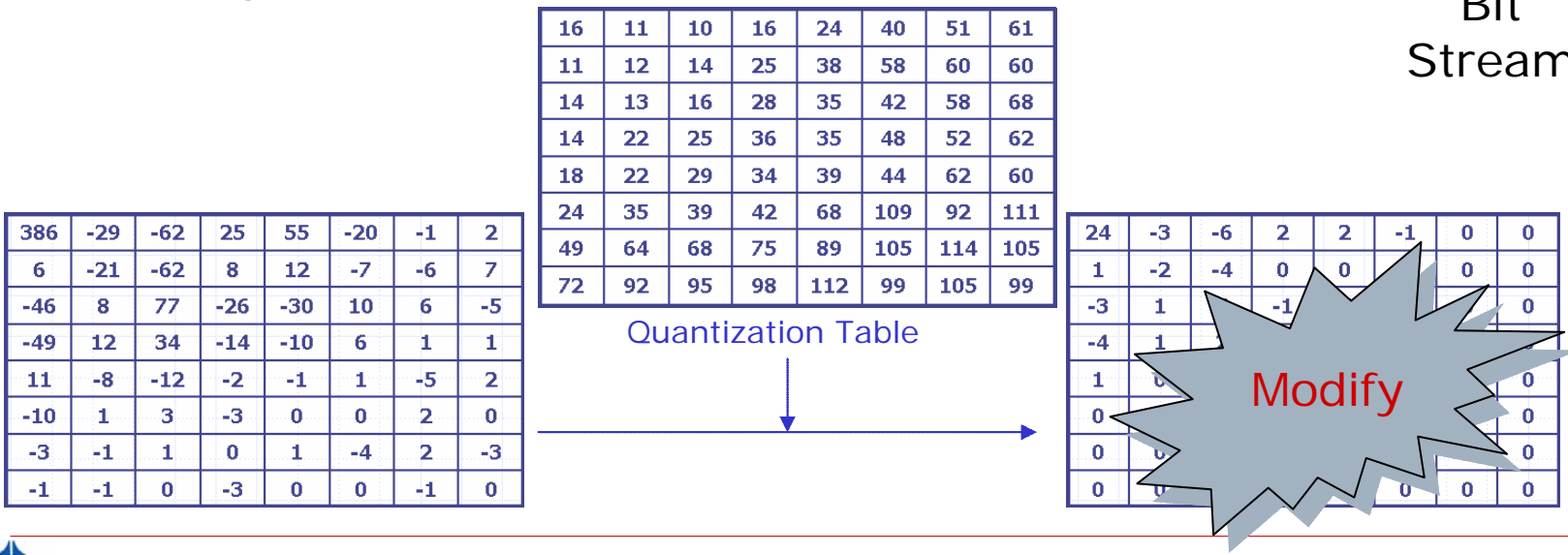
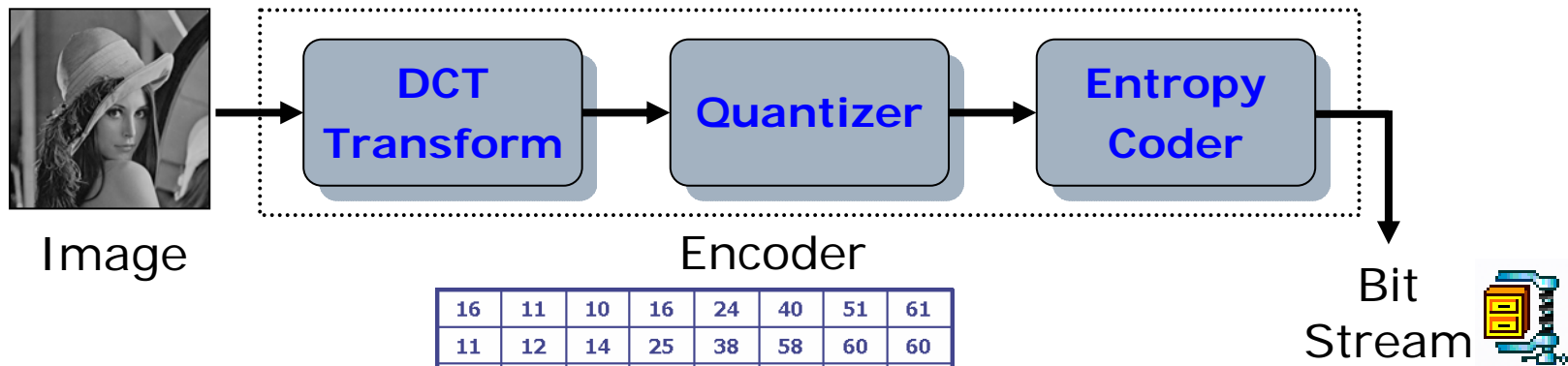
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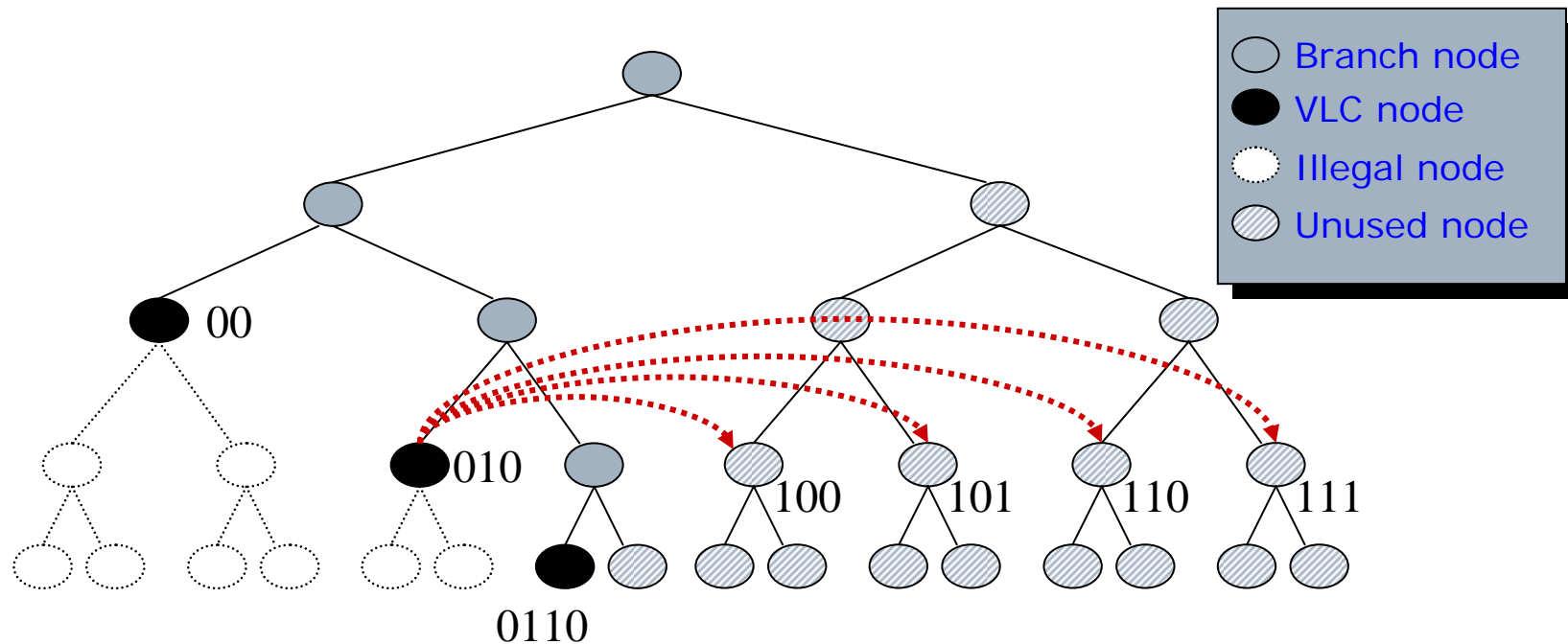
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JPEG Compression



Proposed scheme

- Code tree: $V = \{00, 010, 0110\}$
 - Step 1: Build a code tree of Huffman Coding



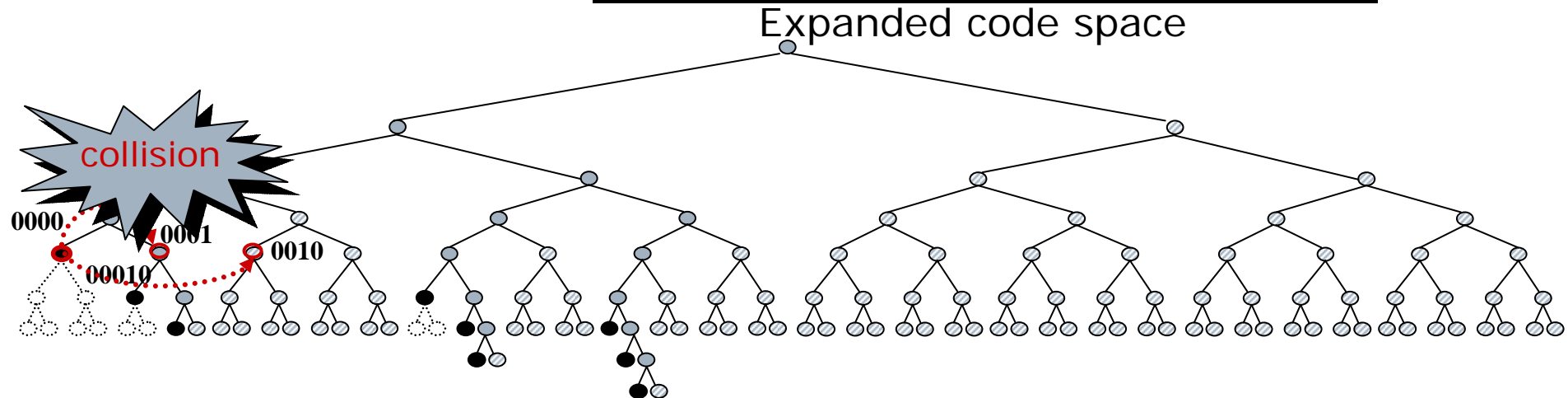
Map a VLC code to an unused codeword for watermarking.

Proposed scheme (cont.)

□ Code tree: $V = \{00, 010, 0110\}$

■ Step 2: Expand the total code space by pairing

1.(00 00)	2.(00 010)	3.(00 0110)
4.(010 00)	5.(010 010)	6.(010 0110)
7.(0110 00)	8.(0110 010)	9.(0110 0110)

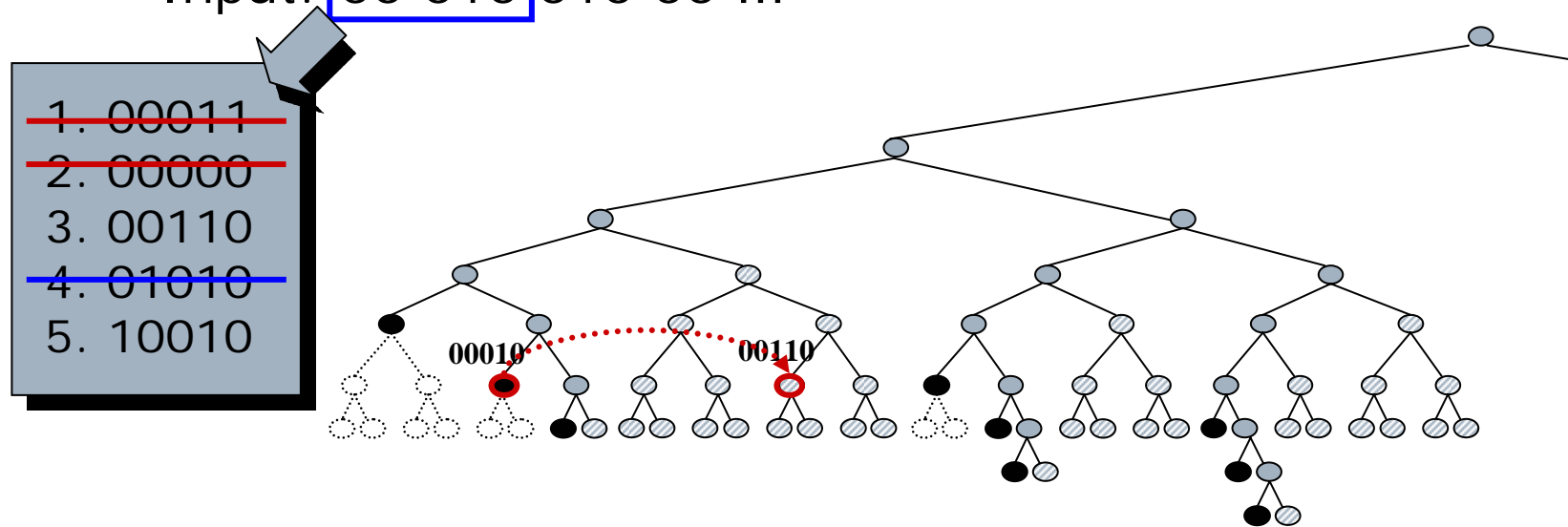


Proposed scheme (cont.)

□ Embedding:

■ Secret: 10...

■ Input: 00 010 010 00 ...

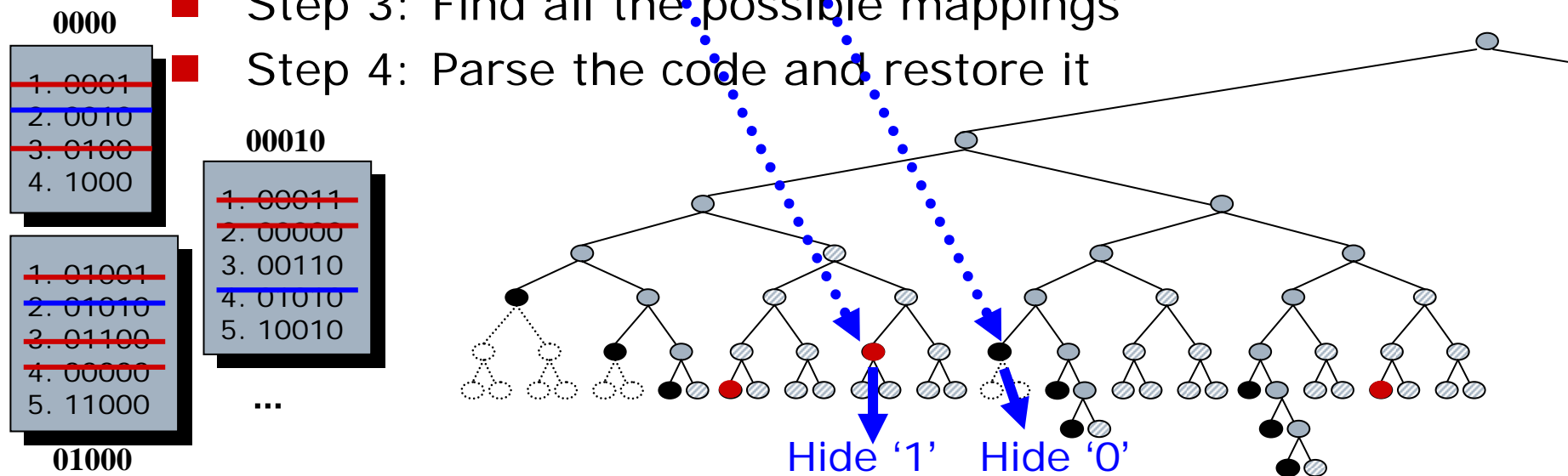


➡ Output: 00 110 010 00 ...

Proposed scheme (cont.)

□ Extracting:

- Received: 0011001000...
- Step 1: Build the original code tree
- Step 2: Build the expanded code tree
- Step 3: Find all the possible mappings
- Step 4: Parse the code and restore it



Experiment

TABLE I
WATERMARKING STATISTICS FOR *LENA.JPG* IMAGE

Total VLCs appearing in image	58
Watermarkable VLCs	31
File size for watermarking (bytes)	46,733
Payload (bits)	2300
Watermark capacity	0.61%

Conclusions

- ❑ Secret data is directly embedded into the compressed *bit stream*.
- ❑ The original compressed bit stream can be *restored*.
- ❑ It is applicable to any entropy-coded compression stream, provided that the total code space is not used.

